

ELIMINATION OF QUANTIFIERS FOR MODULES

BY
WALTER BAUR[†]

ABSTRACT

Every first-order formula in the language of R -modules (R an associative ring) is equivalent relative to the theory of R -modules to a boolean combination of positive primitive formulas and $\forall\exists$ -sentence.

0. Introduction

Let R be an arbitrary associative ring with 1, and let L_R be the language of left R -modules. (For details concerning L_R see, e.g., Eklof-Sabbagh [3]). In the following "module" always means "left R module". Let T_R be the L_R -theory of R -modules. By a positive primitive (p.p.) formula we mean an L_R -formula of the form $\exists \vec{y} \varphi(\vec{x}, \vec{y})$ where $\varphi(\vec{x}, \vec{y})$ is a conjunction of atomic formulas. Formulas without free variables are called sentences.

In this paper we prove the following:

THEOREM. *Every L_R -formula is equivalent relative to T_R to a boolean combination of $\forall\exists$ -sentences and positive primitive formulas.*

This Theorem generalizes in a natural way the corresponding result for abelian groups (Szmielew [5], see also Eklof-Fisher [2]). If we introduce a new n -place predicate symbol R_φ for every p.p. formula with n free variables and if we adjoin to T_R all sentences of the form $\forall \vec{x} (R_\varphi(\vec{x}) \leftrightarrow \varphi(\vec{x}))$, $\varphi(\vec{x})$ p.p., then the Theorem says that this enlarged theory admits elimination of quantifiers modulo certain sentences.

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1. Proof of the Theorem

Sabbagh [4] has shown that every L_R -sentence is equivalent relative to T_R to a boolean combination of $\forall\exists$ -sentences. Therefore it suffices to prove the following weaker statement:

(*) Every L_R -formula is equivalent relative to T_R to a boolean combination of sentences and p.p. formulas.

REMARK. Although a slight modification of our proof would give a direct proof of the Theorem we prefer to prove (*) only, for reasons of simplicity.

TERMINOLOGY. Let $\vec{a} = \langle a_0, \dots, a_{n-1} \rangle$ be an n -tuple of elements from some module M . The set of all p.p. formulas $\varphi(x_0, \dots, x_{n-1})$ such that $M \models \varphi(\vec{a})$ is called the p.p. type of \vec{a} . An isomorphism between submodules A, A' of M is called positive primitive if $M \models \varphi(\vec{a})$ if and only if $M \models \varphi(f(\vec{a}))$ for every p.p. formula $\varphi(\vec{x})$ and every \vec{a} from A . It is easy to see that two n -tuples \vec{a}, \vec{a}' realize the same p.p. type if and only if there exists a p.p. isomorphism f from the submodule generated by the a_i 's onto the submodule generated by the a'_i 's such that $f(a_i) = a'_i$. Whenever we are given n -tuples \vec{a}, \vec{a}' realizing the same p.p. type we will write A for $\sum_{i < n} Ra_i$, A' for $\sum_{i < n} Ra'_i$ and a' for $f(a)$ ($a \in A$). This makes sense since f is unique.

Finally note that all p.p. formulas $\varphi(\vec{x})$ have the following additivity property: $T_R \vdash \forall \vec{x}, \vec{y} (\varphi(\vec{x}) \& \varphi(\vec{y}) \rightarrow \varphi(\vec{x} + \vec{y}))$ and $T_R \vdash \forall \vec{x} (\varphi(\vec{x}) \rightarrow \varphi(\lambda \vec{x}))$ for all $\lambda \in R$.

In this section we will prove (*) modulo the following Lemma whose proof is postponed to the next section:

LEMMA 1. *Let \vec{a}, \vec{a}' be n -tuples of elements from some module M realizing the same p.p. type. Let $\varphi(\vec{x}, y)$ be a p.p. formula, and let $\psi(\vec{x}, y)$ be a conjunction of negations of p.p. formulas. Then $M \models \exists y (\varphi(\vec{a}, y) \& \psi(\vec{a}, y))$ if and only if $M \models \exists y (\varphi(\vec{a}', y) \& \psi(\vec{a}', y))$.*

LEMMA 2. *If two n -tuples \vec{a}, \vec{a}' of elements from some module M realize the same p.p. type then they realize the same elementary type.*

PROOF. Replacing M by some elementary extension we may assume that M is countably saturated. We construct a Karp-isomorphism between the structures $\langle M, \vec{a} \rangle$ and $\langle M, \vec{a}' \rangle$ (see, e.g., Barwise [1]).

Let f be the unique p.p. isomorphism $A \rightarrow A'$ such that $f(a_i) = a'_i$, and let I be the set of all p.p. isomorphisms $g \supseteq f$ between finitely generated submodules of M . By symmetry it suffices to show: if $g \in I$ and $c \in M$ then there exists $h \in I$ such that $g \subseteq h$ and $c \in \text{dom}(h)$.

Put $B = \text{dom}(g)$ and let b_0, \dots, b_{m-1} be a system of generators for B . Let $\Sigma(x_0, \dots, x_{m-1}, y)$ be the set of all conjunctions $\chi(x_0, \dots, x_{m-1}, y)$ of p.p. and negated p.p. formulas such that $M \models \chi(\vec{b}, c)$. By Lemma 1 $M \models \exists y \chi(g(\vec{b}), y)$. Therefore, using saturation, there exists $c' \in M$ such that $\langle g(\vec{b}), c' \rangle$ realizes $\Sigma(\vec{x}, y)$. This implies that $\langle \vec{b}, c \rangle, \langle g(\vec{b}), c' \rangle$ realize the same p.p. type, and hence there exists $h \in I$ such that $h \supseteq g, c \in \text{dom}(h)$ and $h(c) = c'$.

PROOF OF (*). Let $\psi(\vec{x})$ be an arbitrary formula. Let $\Gamma(\vec{x})$ be the set of all formulas $\varphi(\vec{x})$ which are boolean combinations of sentences and p.p. formulas such that $T_R \vdash \varphi(\vec{x}) \rightarrow \psi(\vec{x})$. By compactness it suffices to show that $T = T_R \cup \{\psi(\vec{c})\} \cup \{\neg \varphi(\vec{c}) \mid \varphi(\vec{x}) \in \Gamma(\vec{x})\}$ is inconsistent, where \vec{c} is an n -tuple of new constants. Assume there exists a model $\langle M, \vec{a} \rangle$ of T . Applying Lemma 2 and compactness we obtain a sentence α and a conjunction $\beta(\vec{x})$ of p.p. and negated p.p. formulas such that $M \models \alpha \ \& \ \beta(\vec{a})$ and $T_R \vdash \alpha \ \& \ \beta(\vec{x}) \rightarrow \psi(\vec{x})$. Therefore $\alpha \ \& \ \beta(\vec{x}) \in \Gamma(\vec{x})$ hence $M \models \neg(\alpha \ \& \ \beta(\vec{a}))$, contradiction.

2. Proof of Lemma 1

We will need two auxiliary results.

LEMMA 3. *Let M_0, \dots, M_{k-1} be R -modules and let $\langle c_0, \dots, c_{k-1} \rangle \in M = \bigoplus_{\kappa < k} M_\kappa$. If M' is a submodule of M such that the image of M' under the natural projection $\pi_\kappa: M \rightarrow M_\kappa$ contains more than k elements for every $\kappa < k$, then M' contains an element $\langle b_0, \dots, b_{k-1} \rangle$ such that $b_\kappa \neq c_\kappa$ for all $\kappa < k$.*

The proof is by induction on k . The case $k = 1$ is trivial.

$k > 1$. Case 1. M' infinite

It is easy to see that there exist an infinite subset S of M' and a nonempty subset I of $\{0, \dots, k-1\}$ such that whenever $\vec{b} = \langle b_0, \dots, b_{k-1} \rangle \in S$ and $\vec{b}' = \langle b'_0, \dots, b'_{k-1} \rangle \in S, \vec{b} \neq \vec{b}'$, then $b_\kappa \neq b'_\kappa$ for all $\kappa \in I$ and $b_\kappa = b'_\kappa = 0$ for all $\kappa \notin I$. Applying the induction hypothesis to $N = \bigoplus_{\kappa < k, \kappa \notin I} M_\kappa$ and $N' = \pi(M')$ where $\pi: M \rightarrow N$ is the projection we obtain an element $\vec{b}' = \langle b'_0, \dots, b'_{k-1} \rangle \in M'$ such that $b_\kappa \neq c_\kappa$ for all $\kappa \notin I$. By adding a suitable element from S to \vec{b}' we get an element $\vec{b} \in M'$ satisfying $b_\kappa \neq c_\kappa$ for all $\kappa < k$.

Case 2. M' finite

Let $m = \text{card}(M')$. Assume that for every $\vec{b} = \langle b_0, \dots, b_{k-1} \rangle \in M'$ we have $b_\kappa = c_\kappa$ for some κ .

Put $m_\kappa = \text{card}((\ker \pi_\kappa) \cap M'), l_\kappa = \text{card} \pi_\kappa(M')$. Clearly $m = \sum m_\kappa$ and there-

fore $m_\kappa = m/l_\kappa < m/k$. Since there are at most m_κ elements in M' whose κ -th coordinate is c_κ and since $M' = \bigcup_{\kappa < k} \pi_\kappa^{-1}(c_\kappa)$ by assumption, we obtain

$$m \leq \sum_{\kappa < k} \text{card } \pi_\kappa^{-1}(c_\kappa) \leq \sum_{\kappa < k} m_\kappa < k \frac{m}{k} = m,$$

contradiction.

LEMMA 4. Let \vec{a}, \vec{a}' be n -tuples of elements from some countably saturated module M , and let $\psi(\vec{x}, y)$ be an arbitrary formula. Assume that for all p.p. formulas $\varphi(\vec{x}, y)$, $M \models \exists y(\varphi(\vec{a}, y) \ \& \ \psi(\vec{a}, y))$ if and only if $M \models \exists y(\varphi(\vec{a}', y) \ \& \ \psi(\vec{a}', y))$. If $M \models \exists y\psi(\vec{a}, y)$ then there exist $d, e \in M$ such that

- i) $\langle \vec{a}, d \rangle, \langle \vec{a}', e \rangle$ realize the same p.p. type,
- ii) $M \models \psi(\vec{a}, d)$.

PROOF. For $b \in M$ let Σ_b be the p.p. type of $\langle \vec{a}, b \rangle$. By saturation the set $\{\Sigma_b \mid M \models \psi(\vec{a}, b)\}$ contains a maximal element $\Sigma_d, d \in M$. Using the hypothesis and saturation again we obtain an element $e \in M$ such that $M \models \varphi(\vec{a}', e) \ \& \ \psi(\vec{a}', e)$ for all $\varphi(\vec{x}, y) \in \Sigma_d$. Let $\varphi(\vec{x}, y)$ be a p.p. formula such that $M \models \varphi(\vec{a}', e)$. Arguing as above we see that $\Sigma_d(\vec{a}, y) \cup \{\varphi(\vec{a}, y)\} \cup \{\psi(\vec{a}, y)\}$ is consistent. Hence $\varphi(\vec{x}, y) \in \Sigma_d$, by maximality. Therefore $\langle \vec{a}, d \rangle$ and $\langle \vec{a}', e \rangle$ realize the same p.p. type.

PROOF OF LEMMA 1. We proceed by induction on the number k of conjuncts in $\psi(\vec{x}, y)$. The case $k = 0$ is just the hypothesis. Let $k > 0$ and assume Lemma 1 to be true for all $k' < k$ and all n . By symmetry it suffices to show that

$$M \models \exists y(\varphi(\vec{a}, y) \ \& \ \psi(\vec{a}, y)) \text{ implies } M \models \exists y(\varphi(\vec{a}', y) \ \& \ \psi(\vec{a}', y)).$$

Therefore let $b_1 \in M$ such that

$$(1) \quad M \models \varphi(\vec{a}, b_1) \ \& \ \psi(\vec{a}, b_1).$$

Furthermore we may assume that M is countably saturated.

Write $\psi(\vec{x}, y)$ as $\bigwedge_{\kappa < k} \neg \psi_\kappa(\vec{x}, y)$, $\psi_\kappa(\vec{x}, y)$ p.p.

Put

$$N = \{d \in M \mid M \models \varphi(\vec{0}, d)\},$$

$$P_\kappa = M^{n+1} / \{(\vec{c}, d) \mid M \models \psi_\kappa(\vec{c}, d)\},$$

and let $q_\kappa : M^{n+1} \rightarrow P_\kappa$ be the natural projection. Put $P = \bigoplus_{\kappa < k} P_\kappa$ and let Q be the image of N under the map $M \rightarrow P$ defined by

$$d \mapsto \langle q_0(\vec{0}, d), \dots, q_{k-1}(\vec{0}, d) \rangle.$$

Case 1. The image of Q under the projection $P \rightarrow P_\kappa$ contains more than k elements for each $\kappa < k$.

Since \vec{a}, \vec{a}' realize the same p.p. type there exists $b \in M$ such that $M \models \varphi(\vec{a}', b)$. Applying Lemma 3 ($M_\kappa = P_\kappa, M' = Q$) we find $d \in N$ such that $q_\kappa(\vec{0}, d) \neq q_\kappa(\vec{a}', b)$ for all $\kappa < k$, in other words $M \models \bigwedge_{\kappa < k} \neg \psi_\kappa(\vec{a}', b - d)$. Since $M \models \varphi(\vec{a}', b - d)$ by additivity we conclude $M \models \exists y (\varphi(\vec{a}', y) \ \& \ \psi(\vec{a}', y))$ as desired.

Case 2. There exist $\lambda < k, h \leq k$ such that the image J of Q under the projection $P \rightarrow P_\lambda$ contains exactly h elements.

Let J' be the set consisting of those elements $p_\lambda \in J$ such that there exists $d \in A \cap N$ with $q_\lambda(\vec{0}, d) = p_\lambda$. Let $h' (\leq h)$ be the cardinality of J' and choose $d_0, \dots, d_{h'-1} \in A \cap N$ such that $\{q_\lambda(\vec{0}, d_i) \mid i < h'\} = J'$.

Case 2.1. $h - h' = 0$, i.e. $J' = J$. We may assume that there exists an element $b \in A$ such that $M \models \varphi(\vec{a}, b)$. (If A does not contain such an element we can adjoin one by applying Lemma 4 with $\varphi(\vec{x}, y)$ as $\psi(\vec{x}, y)$). The element $b_1 - b$ lies in N , by the additivity of $\varphi(\vec{x}, y)$. Since $J' = J$ there exists $d \in A \cap N$ such that $q_\lambda(\vec{0}, d) = q_\lambda(\vec{0}, b_1 - b)$. Therefore if $a_n = b + d$ then

$$(2) \quad M \models \psi_\lambda(\vec{0}, b_1 - a_n).$$

Let $\chi(x_0, \dots, x_n, y)$ be the formula

$$\varphi(x_0, \dots, x_{n-1}, y) \ \& \ \psi_\lambda(\vec{0}, y - x_n) \ \& \ \bigwedge_{\substack{\kappa < k \\ \kappa \neq \lambda}} \neg \psi_\kappa(\vec{x}, y).$$

By (1) and (2), $M \models \chi(\vec{a}, a_n, b_1)$ and therefore $M \models \exists y \chi(\vec{a}, a_n, y)$. Since $a_n \in A$, the $(n + 1)$ -tuples $\langle a_0, \dots, a_n \rangle, \langle a'_0, \dots, a'_n \rangle$ realize the same p.p. type. (Recall that a'_v is the image of a_v under the p.p. isomorphism $a_v \mapsto a'_v$ ($v < n$)). Since $\chi(\vec{x}, y)$ contains only $k - 1$ conjuncts being negations of p.p. formulas we conclude $M \models \exists y \chi(\vec{a}', a'_n, y)$, by induction hypothesis. Let $b_2 \in M$ such that $M \models \chi(\vec{a}', a'_n, b_2)$. This means

$$M \models \varphi(\vec{a}', b_2) \ \& \ \bigwedge_{\substack{\kappa < k \\ \kappa \neq \lambda}} \neg \psi_\kappa(\vec{a}', b_2) \text{ and}$$

$$(3) \quad M \models \psi_\lambda(\vec{0}, b_2 - a'_n).$$

It remains to show that $M \models \neg \psi_\lambda(\vec{a}', b_2)$. If $M \models \psi_\lambda(\vec{a}', b_2)$ then $M \models \psi_\lambda(\vec{a}', a'_n)$ by (3) and additivity, hence $M \models \psi_\lambda(\vec{a}, a_n)$, and therefore $M \models \psi_\lambda(\vec{a}, b_1)$ by (2), contradicting (1).

Case 2.2. $h - h' > 0$. We reduce this case to the former by successively adjoining elements from M to A and A' . Clearly it suffices to prove the following: There exist $d, e \in N$ such that

- i) $\langle \vec{a}, d \rangle, \langle \vec{a}', e \rangle$ realize the same p.p. type,
- ii) $q_\lambda(\langle \vec{0}, d \rangle) \in J - J'$.

Let $\psi'(u_0, \dots, u_{h'-1}, y)$ be the formula

$$\varphi(\vec{0}, y) \ \& \ \bigwedge_{i < h'} \neg \psi_\lambda(\vec{0}, y - u_i).$$

Clearly $M \models \psi'(d_0, \dots, d_{h'-1}, d)$ if and only if $d \in N$ and $q_\lambda(\langle \vec{0}, d \rangle) \notin J'$. Since the d_i 's lie in A , the $(n + h')$ -tuples $\langle a_0, \dots, a_{n-1}, d_0, \dots, d_{h'-1} \rangle, \langle a'_0, \dots, a'_{n-1}, d'_0, \dots, d'_{h'-1} \rangle$ realize the same p.p. type. Finally note that $\psi'(\vec{u}, y)$ contains only $h' < h \leq k$ conjuncts being negations of p.p. formulas. Combining the last two facts with the induction hypothesis and Lemma 4 (with $\psi'(\vec{u}, y)$ as $\psi(\vec{x}, y)$) we obtain elements $d, e \in N$ satisfying (i) and (ii). This concludes the proof of Lemma 1.

3. Abelian groups

When dealing with abelian groups we can replace p.p. formulas by formulas expressing divisibility. Call a formula $\varphi(x_0, \dots, x_{n-1})$ a d -formula if it is of the form $\sum_{\nu < n} k_\nu x_\nu = 0$ or of the form $\exists y (p^k y = \sum_{\nu < n} k_\nu x_\nu)$ for some prime p , natural number k and integers k_ν . Combining the elementary Lemma 4.3 from [2] with the fact that an embedding between abelian groups preserves p.p. formulas if and only if it preserves d -formulas it is easy to see that two n -tuples \vec{a}, \vec{a}' from some abelian group M realize the same p.p. type if and only if they realize the same d -type. Applying Lemma 2 and compactness as in the proof of (*) we obtain the following weak form of Szmielew's elimination-of-quantifier-result:

Every formula in the language of abelian groups is equivalent relative to the theory of abelian groups to a boolean combination of $\forall \exists$ -sentences and divisibility formulas.

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UNIVERSITY OF ZURICH
FREIESTRASSE 36, SWITZERLAND